



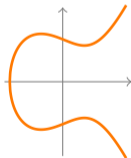
Return of ECC dummy point additions: SPA on efficient P-256 implementation

Andy Russon

3 June 2021

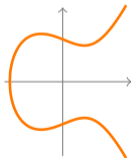


- Secure cryptography: avoid secret dependent behavior in execution



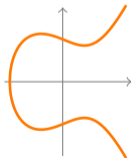
¹Gueron and Krasnov, "Fast Prime Field Elliptic Curve Cryptography with 256 Bit Prime", 2013

- Secure cryptography: avoid secret dependent behavior in execution
- Dummy operations might achieve this goal



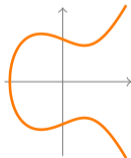
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- Example of elliptic curve P-256:



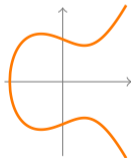
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 - Popular curve, notably for digital signatures (with ECDSA)



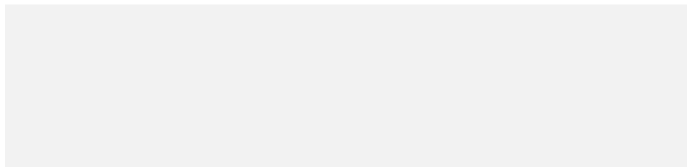
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- Secure cryptography: avoid secret dependent behavior in execution
- Dummy operations might achieve this goal
- Example of elliptic curve P-256:
 - Popular curve, notably for digital signatures (with ECDSA)
 - Efficient implementation in assembly available (OpenSSL and its forks)¹



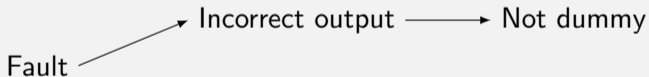
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- Previously at SSTIC²
 - Fault injection to detect dummy operation



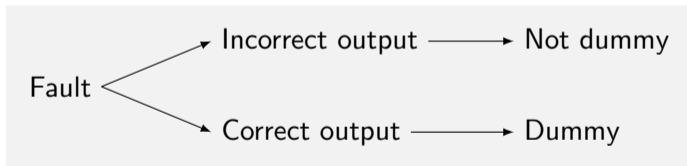
²“Exploiting dummy codes in Elliptic Curve Cryptography implementations”, SSTIC 2020

- Previously at SSTIC²
 - Fault injection to detect dummy operation



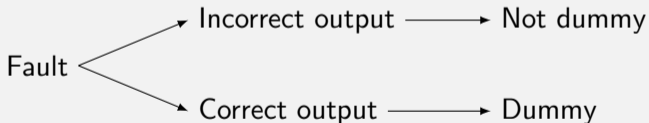
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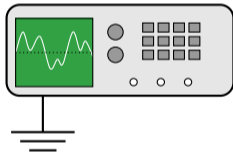
²“Exploiting dummy codes in Elliptic Curve Cryptography implementations”, SSTIC 2020

- Previously at SSTIC²
 - Fault injection to detect dummy operation
 - Not easy (material, timing, etc)
 - Produces many invalid outputs: hard to pass unnoticed



²“Exploiting dummy codes in Elliptic Curve Cryptography implementations”, SSTIC 2020

- Goal: recover an ECDSA private key
 - On OpenSSL with the P-256 assembly optimized implementation
 - With power analysis to detect dummy operations

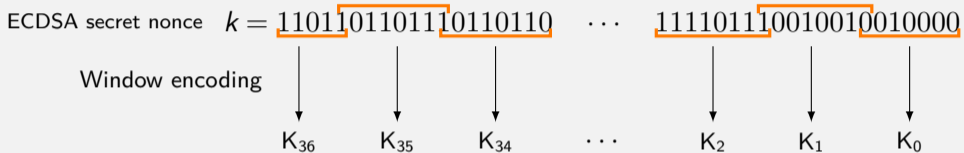


- 1 Overview of the target algorithm
- 2 Dummy addition and power analysis
- 3 Dummy addition and zeros
- 4 Conclusion and mitigations

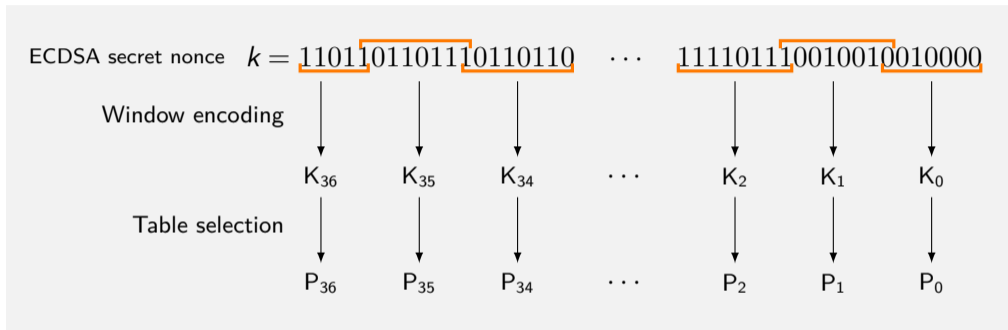
ECDSA secret nonce $k = 1101101101110110110 \dots 111101110010010010000$

ECDSA secret nonce $k =$ 11011011011110110110 ... 1111011110010010010000

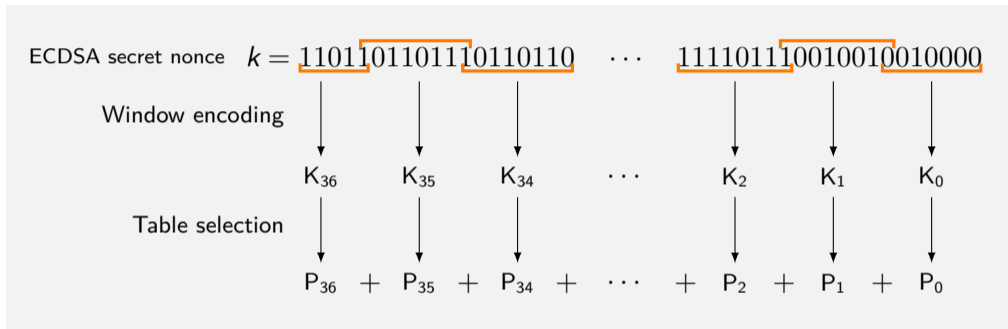
Assembly optimized P-256 implementation



Assembly optimized P-256 implementation

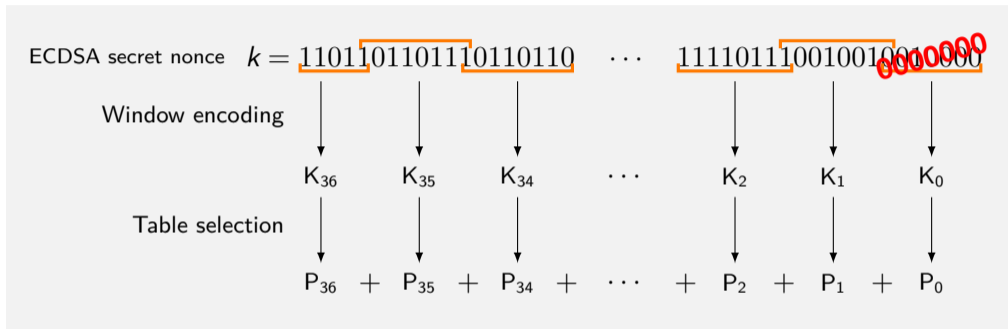


Assembly optimized P-256 implementation



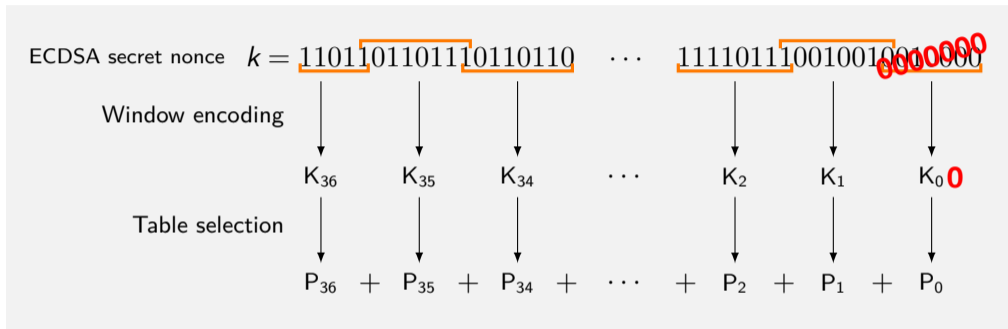
```
function ecp_nistz_point_add_affine
```

Assembly optimized P-256 implementation



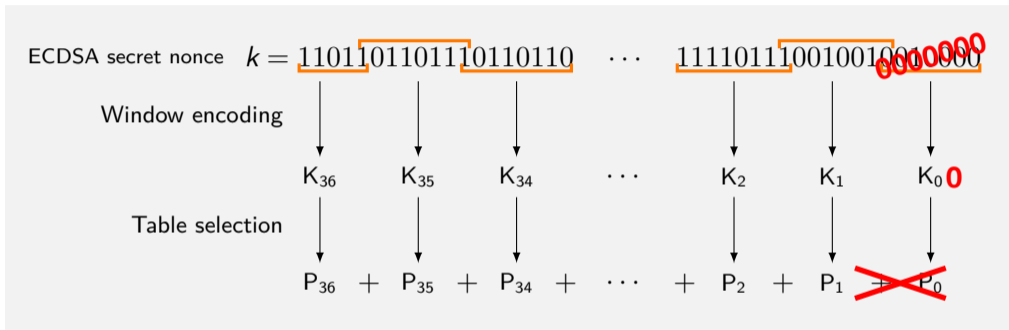
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Assembly optimized P-256 implementation



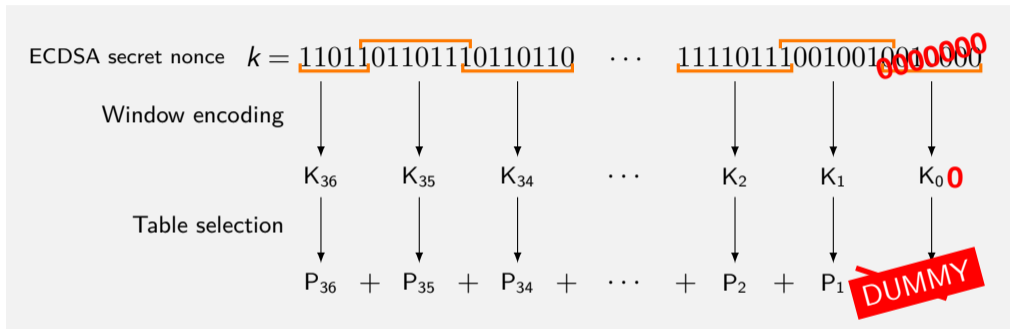
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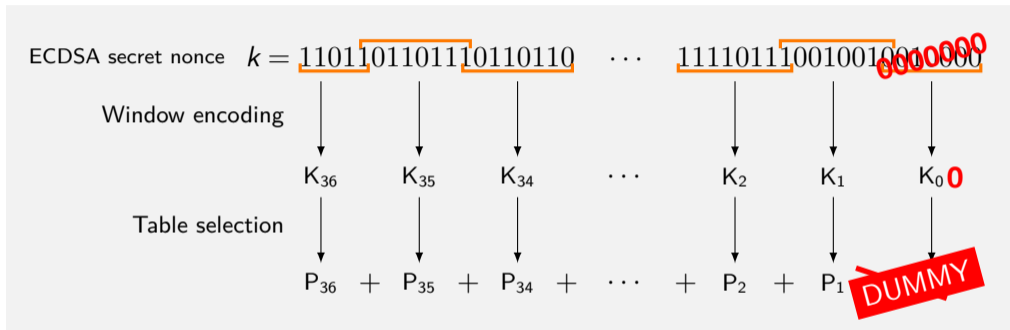
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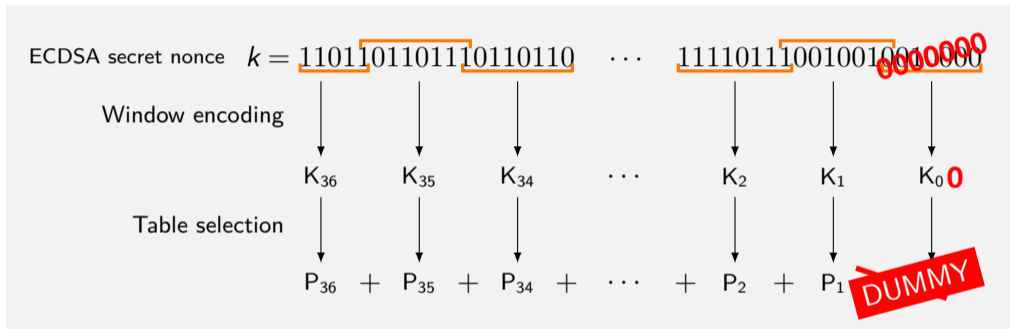


- Always 36 execution function `ecp_nistz_point_add_affine`

Assembly optimized P-256 implementation



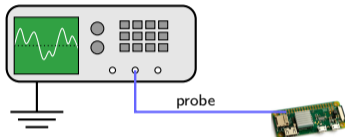
- Always 36 execution function `ecp_nistz_point_add_affine`
- Null window \Rightarrow dummy point addition



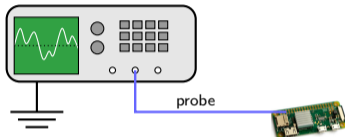
- Always 36 execution function `ecp_nistz_point_add_affine`
- Null window \Rightarrow dummy point addition
- First window: K_0 is 7 lowest bits of k

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- 1 Capture power consumption of an ECDSA signature generation

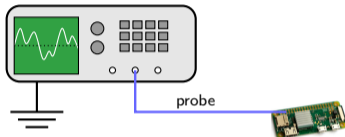


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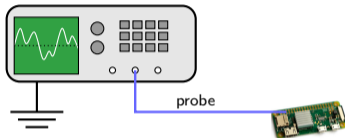
- 2 Keep the signature if there is a lower power consumption during the **first** point addition

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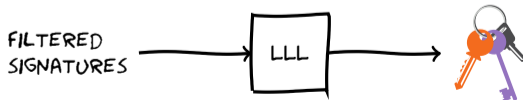


- 2 Keep the signature if there is a lower power consumption during the **first** point addition
- 3 Repeat until enough signatures are collected (around 37)

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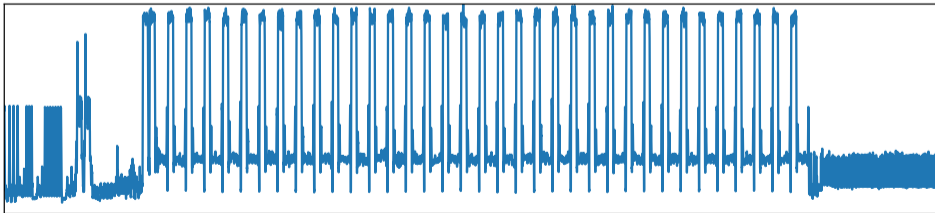


- 2 Keep the signature if there is a lower power consumption during the **first** point addition
- 3 Repeat until enough signatures are collected (around 37)
- 4 Recover the private key from the signatures

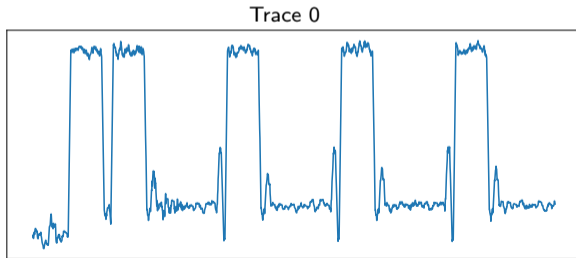


- TracerGrind: a Valgrind plugin to monitor processes
- Simulation of power traces:
 - Capture memory trace of a signature generation
 - Apply Hamming weight model

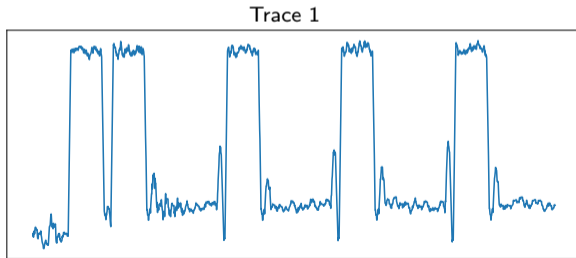
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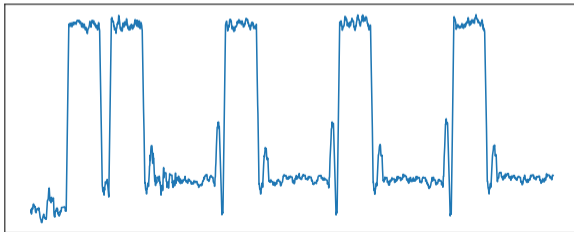


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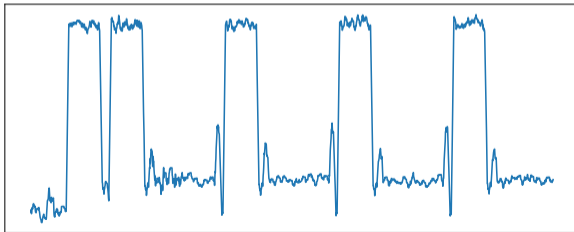
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Trace 2



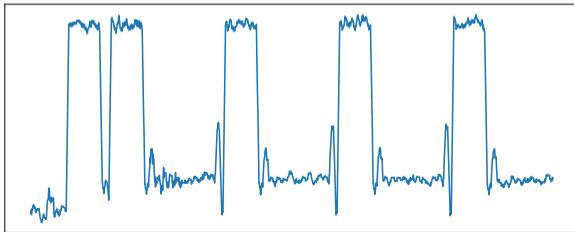
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Trace 3



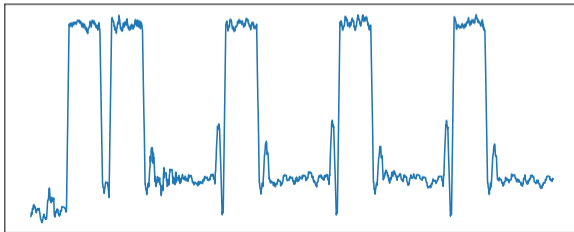
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Trace 4



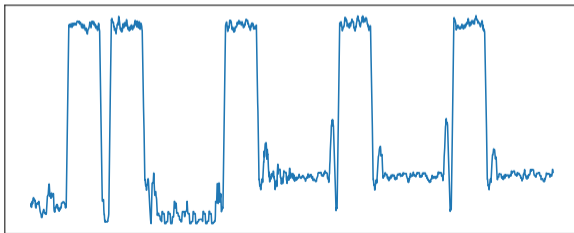
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Trace 5

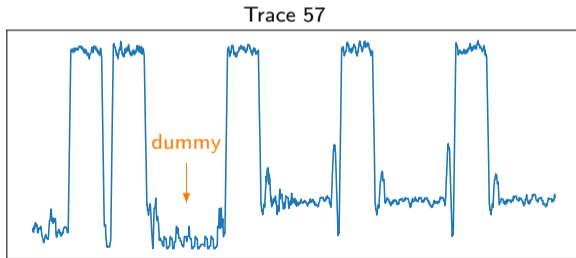


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Trace 57



- TracerGrind: a Valgrind plugin to monitor processes
- Simulation of power traces:
 - Capture memory trace of a signature generation
 - Apply Hamming weight model



How many signatures we need?



- Number of signatures:
 - ▶ 7 lowest bits of k null: 1 out of 128 signatures on average
 - ▶ Key recovery: 37 signatures (or more) if 7 bits of nonce known
 - ▶ Total: $37 \times 128 = 4738$

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 - ▶ 7 lowest bits of k null: 1 out of 128 signatures on average
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 - ▶ Total: $37 \times 128 = 4738$
- Clustering the traces:
 - ▶ GaussianMixture class in Scikit-learn Python library
 - ▶ Classify the traces in **two** components

```
from sklearn import mixture
gm = mixture.GaussianMixture(n_components=2)
results = gm.fit_predict(traces)
```


- Trace generation:

```
python3 simulation.py -b /usr/local/bin/openssl --privkey privkey.pem  
--working-dir test -n 6000 --bounds 1844200 1844500
```

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- Analysis:

```
python3 analysis.py --pubkey pubkey.pem --working-dir test -n 6000
```

- Trace generation:

```
python3 simulation.py -b /usr/local/bin/openssl --privkey privkey.pem  
--working-dir test -n 6000 --bounds 1844200 1844500
```

- Analysis:

```
python3 analysis.py --pubkey pubkey.pem --working-dir test -n 6000
```

- Results:

```
Loading 6000 traces...  
Clustering of the traces...  
Size of clusters: 5950, 50  
Lattice attack: recovering of the key with 37 signatures...  
Lattice attack: recovering of the key with 38 signatures...  
SUCCESS!  
The private key is deb867c4cf52002d7d8cc2e440d16c8c75c7844fdb142ea8c1a1f48b411b0251
```

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`ecp_nistz_point_add_affine:`

Input: $P = (X_1, Y_1, Z_1)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

```
in1infy  $\leftarrow (Z_1 == 0)$ 
in2infy  $\leftarrow (x_2, y_2) == (0, 0)$ 
t0  $\leftarrow Z_1^2$ 
t1  $\leftarrow x_2 \cdot t_0$ 
t2  $\leftarrow t_1 - X_1$ 
t3  $\leftarrow t_0 \cdot Z_1$ 
Z3  $\leftarrow t_2 \cdot Z_1$ 
t3  $\leftarrow t_3 \cdot y_2$ 
t4  $\leftarrow t_3 - Y_1$ 
t5  $\leftarrow t_2^2$ 
t6  $\leftarrow t_4^2$ 
t7  $\leftarrow t_5 \cdot t_2$ 
t1  $\leftarrow X_1 \cdot t_5$ 
t5  $\leftarrow 2 \cdot t_1$ 
X3  $\leftarrow t_6 - t_5$ 
X3  $\leftarrow X_3 - t_7$ 
t2  $\leftarrow t_1 - X_3$ 
t3  $\leftarrow Y_1 \cdot t_7$ 
t2  $\leftarrow t_2 \cdot t_4$ 
Y3  $\leftarrow t_2 - t_3$ 
if in1infy then
     $(X_3, Y_3, Z_3) \leftarrow (x_2, y_2, 1)$ 
if in2infy then
     $(X_3, Y_3, Z_3) \leftarrow (X_1, Y_1, Z_1)$ 
return  $(X_3, Y_3, Z_3)$ 
```

Handling of the dummy point addition



- If $K_0 = 0$, then P is initialize with $(0, 0, 0)$
- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

```
in1infty ← (Z1 == 0)
in2infty ← (x2, y2) == (0, 0)
t0 ← Z12
t1 ← x2 · t0
t2 ← t1 - X1
t3 ← t0 · Z1
Z3 ← t2 · Z1
t3 ← t3 · y2
t4 ← t3 - Y1
t5 ← t22
t6 ← t42
t7 ← t5 · t2
t1 ← X1 · t5
t5 ← 2 · t1
X3 ← t6 - t5
X3 ← X3 - t7
t2 ← t1 - X3
t3 ← Y1 · t7
t2 ← t2 · t4
Y3 ← t2 - t3
if in1infty then
    (X3, Y3, Z3) ← (x2, y2, 1)
if in2infty then
    (X3, Y3, Z3) ← (X1, Y1, Z1)
return (X3, Y3, Z3)
```

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```

in1infy ← true
in2infy ← (x2, y2) == (0, 0)
t0 ← Z12
t1 ← x2 · t0
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t3 ← t0 · Z1
Z3 ← t2 · Z1
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if in1infy then
    (X3, Y3, Z3) ← (x2, y2, 1)
if in2infy then
    (X3, Y3, Z3) ← (X1, Y1, Z1)
return (X3, Y3, Z3)

```

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- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

<code>in1infty</code> \leftarrow true	$t_5 \leftarrow 2 \cdot t_1$
<code>in2infty</code> $\leftarrow (x_2, y_2) == (0, 0)$	$X_3 \leftarrow t_6 - t_5$
$t_0 \leftarrow 0^2$	$X_3 \leftarrow X_3 - t_7$
$t_1 \leftarrow x_2 \cdot t_0$	$t_2 \leftarrow t_1 - X_3$
$t_2 \leftarrow t_1 - X_1$	$t_3 \leftarrow Y_1 \cdot t_7$
$t_3 \leftarrow t_0 \cdot Z_1$	$t_2 \leftarrow t_2 \cdot t_4$
$Z_3 \leftarrow t_2 \cdot Z_1$	$Y_3 \leftarrow t_2 - t_3$
$t_3 \leftarrow t_3 \cdot y_2$	if <code>in1infty</code> then
$t_4 \leftarrow t_3 - Y_1$	$(X_3, Y_3, Z_3) \leftarrow (x_2, y_2, 1)$
$t_5 \leftarrow t_2^2$	if <code>in2infty</code> then
$t_6 \leftarrow t_4^2$	$(X_3, Y_3, Z_3) \leftarrow (X_1, Y_1, Z_1)$
$t_7 \leftarrow t_5 \cdot t_2$	return (X_3, Y_3, Z_3)
$t_1 \leftarrow X_1 \cdot t_5$	

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```

in1infty ← true
in2infty ← (x2, y2) == (0, 0)
t0 ← 02
t1 ← x2 · 0
t2 ← t1 - X1
t3 ← t0 · Z1
Z3 ← t2 · Z1
t3 ← t3 · y2
t4 ← t3 - Y1
t5 ← t22
t6 ← t42
t7 ← t5 · t2
t1 ← X1 · t5
t5 ← 2 · t1
X3 ← t6 - t5
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t2 ← t1 - X3
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t2 ← t2 · t4
Y3 ← t2 - t3
if in1infty then
    (X3, Y3, Z3) ← (x2, y2, 1)
if in2infty then
    (X3, Y3, Z3) ← (X1, Y1, Z1)
return (X3, Y3, Z3)

```

Handling of the dummy point addition

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- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

```

in1infty ← true
in2infty ← (x2, y2) == (0, 0)
t0 ← 02
t1 ← x2 · 0
t2 ← 0 - 0
t3 ← t0 · Z1
Z3 ← t2 · Z1
t3 ← t3 · y2
t4 ← t3 - Y1
t5 ← t22
t6 ← t42
t7 ← t5 · t2
t1 ← X1 · t5
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if in1infty then
    (X3, Y3, Z3) ← (x2, y2, 1)
if in2infty then
    (X3, Y3, Z3) ← (X1, Y1, Z1)
return (X3, Y3, Z3)

```

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- If $K_0 = 0$, then P is initialize with $(0, 0, 0)$
- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

<code>in1infty</code> \leftarrow true	<code>t₅</code> $\leftarrow 2 \cdot t_1$
<code>in2infty</code> $\leftarrow (x_2, y_2) == (0, 0)$	<code>X₃</code> $\leftarrow t_6 - t_5$
<code>t₀</code> $\leftarrow 0^2$	<code>X₃</code> $\leftarrow X_3 - t_7$
<code>t₁</code> $\leftarrow x_2 \cdot 0$	<code>t₂</code> $\leftarrow t_1 - X_3$
<code>t₂</code> $\leftarrow 0 - 0$	<code>t₃</code> $\leftarrow Y_1 \cdot t_7$
<code>t₃</code> $\leftarrow 0 \cdot 0$	<code>t₂</code> $\leftarrow t_2 \cdot t_4$
<code>Z₃</code> $\leftarrow t_2 \cdot Z_1$	<code>Y₃</code> $\leftarrow t_2 - t_3$
<code>t₃</code> $\leftarrow t_3 \cdot y_2$	if <code>in1infty</code> then
<code>t₄</code> $\leftarrow t_3 - Y_1$	$(X_3, Y_3, Z_3) \leftarrow (x_2, y_2, 1)$
<code>t₅</code> $\leftarrow t_2^2$	if <code>in2infty</code> then
<code>t₆</code> $\leftarrow t_4^2$	$(X_3, Y_3, Z_3) \leftarrow (X_1, Y_1, Z_1)$
<code>t₇</code> $\leftarrow t_5 \cdot t_2$	return (X_3, Y_3, Z_3)
<code>t₁</code> $\leftarrow X_1 \cdot t_5$	

Handling of the dummy point addition

- If $K_0 = 0$, then P is initialize with $(0, 0, 0)$
- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

<code>in1infty</code> \leftarrow true	<code>t₅</code> $\leftarrow 2 \cdot t_1$
<code>in2infty</code> $\leftarrow (x_2, y_2) == (0, 0)$	<code>X₃</code> $\leftarrow t_6 - t_5$
<code>t₀</code> $\leftarrow 0^2$	<code>X₃</code> $\leftarrow X_3 - t_7$
<code>t₁</code> $\leftarrow x_2 \cdot 0$	<code>t₂</code> $\leftarrow t_1 - X_3$
<code>t₂</code> $\leftarrow 0 - 0$	<code>t₃</code> $\leftarrow Y_1 \cdot t_7$
<code>t₃</code> $\leftarrow 0 \cdot 0$	<code>t₂</code> $\leftarrow t_2 \cdot t_4$
<code>Z₃</code> $\leftarrow 0 \cdot 0$	<code>Y₃</code> $\leftarrow t_2 - t_3$
<code>t₃</code> $\leftarrow t_3 \cdot y_2$	if <code>in1infty</code> then
<code>t₄</code> $\leftarrow t_3 - Y_1$	$(X_3, Y_3, Z_3) \leftarrow (x_2, y_2, 1)$
<code>t₅</code> $\leftarrow t_2^2$	if <code>in2infty</code> then
<code>t₆</code> $\leftarrow t_4^2$	$(X_3, Y_3, Z_3) \leftarrow (X_1, Y_1, Z_1)$
<code>t₇</code> $\leftarrow t_5 \cdot t_2$	return (X_3, Y_3, Z_3)
<code>t₁</code> $\leftarrow X_1 \cdot t_5$	

Handling of the dummy point addition

- If $K_0 = 0$, then P is initialize with $(0, 0, 0)$
- Consequence in `ecp_nistz_point_add_affine`:


Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

<code>in1infty</code> \leftarrow true	<code>t₅</code> \leftarrow $2 \cdot 0$
<code>in2infty</code> \leftarrow $(x_2, y_2) == (0, 0)$	<code>X₃</code> \leftarrow $0 - 0$
<code>t₀</code> \leftarrow 0^2	<code>X₃</code> \leftarrow $0 - 0$
<code>t₁</code> \leftarrow $x_2 \cdot 0$	<code>t₂</code> \leftarrow $0 - 0$
<code>t₂</code> \leftarrow $0 - 0$	<code>t₃</code> \leftarrow $0 \cdot 0$
<code>t₃</code> \leftarrow $0 \cdot 0$	<code>t₂</code> \leftarrow $0 \cdot 0$
<code>Z₃</code> \leftarrow $0 \cdot 0$	<code>Y₃</code> \leftarrow $0 - 0$
<code>t₃</code> \leftarrow $0 \cdot y_2$	if <code>in1infty</code> then
<code>t₄</code> \leftarrow $0 - 0$	$(X_3, Y_3, Z_3) \leftarrow (x_2, y_2, 1)$
<code>t₅</code> \leftarrow 0^2	if <code>in2infty</code> then
<code>t₆</code> \leftarrow 0^2	$(X_3, Y_3, Z_3) \leftarrow (X_1, Y_1, Z_1)$
<code>t₇</code> \leftarrow $0 \cdot 0$	return (X_3, Y_3, Z_3)
<code>t₁</code> \leftarrow $0 \cdot 0$	

Handling of the dummy point addition

- If $K_0 = 0$, then P is initialize with $(0, 0, 0)$
- Consequence in `ecp_nistz_point_add_affine`:

Input: $P = (0, 0, 0)$, $Q = (x_2, y_2)$ **Output:** $P + Q = (X_3, Y_3, Z_3)$

<pre> in1infty ← true in2infty ← (x2, y2) == (0, 0) t0 ← 0² t1 ← x2 · 0 t2 ← 0 - 0 t3 ← 0 · 0 Z3 ← 0 · 0 t3 ← 0 · y2 t4 ← 0 - 0 t5 ← 0² t6 ← 0² t7 ← 0 · 0 t1 ← 0 · 0 </pre>	<pre> t5 ← 2 · 0 X3 ← 0 - 0 X3 ← 0 - 0 t2 ← 0 - 0 t3 ← 0 · 0 t2 ← 0 · 0 Y3 ← 0 - 0 </pre>	<p>ignore previous calculation</p> 
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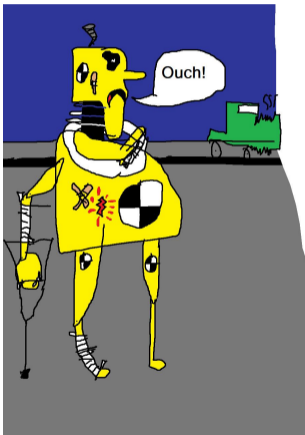
```

if in1infty then
    (X3, Y3, Z3) ← (x2, y2, 1)
        
```

```

if in2infty then
    (X3, Y3, Z3) ← (X1, Y1, Z1)
return (X3, Y3, Z3)
        
```

- 1 Overview of the target algorithm
- 2 Dummy addition and power analysis
- 3 Dummy addition and zeros
- 4 Conclusion and mitigations



- SPA on efficient implementation of elliptic curve P-256
- Less than 5000 traces to recover a private key
- Mitigations:
 - Dummy additions with random values instead of zero
 - Use regular algorithms that do not need dummy additions
- Proof of concept and tools available as a Docker:
<https://github.com/orangecertcc/ecdummyrpa>